

Fully Efficient Nonnegative Matrix Factorization

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This presentation is some rephrasing of the 5th section of the following paper :

- ▶ Arora, S., Ge, R., Kannan, R., and Moitra, A., " *Computing a nonnegative matrix factorization, provably,*" In ACM Symposium on Theory of Computing (STOC), 2012.

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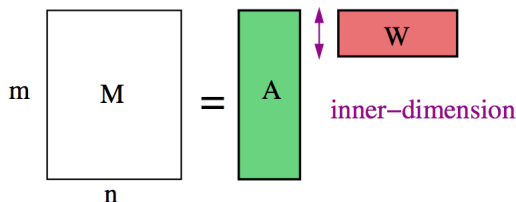
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Problem Definition

- ▶ Noise Free Case : Given a non-negative $m \times n$ matrix \mathbf{M} , find $m \times r$ matrix \mathbf{A} and $r \times n$ matrix \mathbf{W} such that $\mathbf{M} = \mathbf{A}\mathbf{W}$.
- ▶ Small Perturbation Case : Suppose $\mathbf{M} = \mathbf{A}\mathbf{W}$, but instead of \mathbf{M} , only a noisy version \mathbf{M}' is given.
 - ▶ Assumption : $\|\mathbf{M}_i - \mathbf{M}'_i\|_1 \leq \epsilon$.
 - ▶ Goal : Find \mathbf{A}' and \mathbf{W}' such that $\|\mathbf{M}'_j - (\mathbf{A}'\mathbf{W}')_j\|_1 \leq \mathcal{O}(\epsilon)$
- ▶ WLoG : Assume that \mathbf{M} and \mathbf{W} are row normalized.



Picture Courtesy : A. Moitra, "Computing a Nonnegative Matrix Factorization – Provably," Dec. 2011.

Challenges

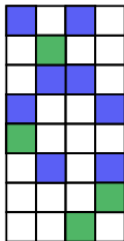
- ▶ The problem can be solved efficiently using SVD, if \mathbf{A} and \mathbf{W} are not required to be nonnegative.
- ▶ It has been recently shown by Vavasis that the problem is \mathcal{NP} -hard.
- ▶ Under the mild separability assumption of \mathbf{A} , Arora has shown that a very efficient algorithm exists to solve the problem.

Separability Assumption

Definition

Matrix \mathbf{A} is a separable matrix, if for each column i , there exists some row $f(i)$ for which the i^{th} element is the only non-zero element in that row. These row indices are called *anchor rows*.

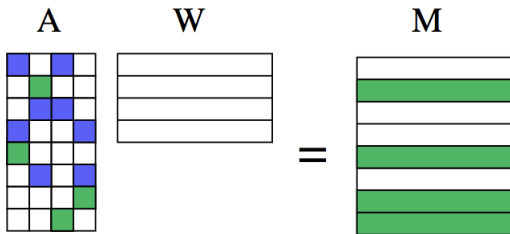
\mathbf{A}



Picture Courtesy : A. Moitra, "Computing a Nonnegative Matrix Factorization – Provably," Dec. 2011.

Why Separability is important?

- ▶ After a proper permutation of its rows, \mathbf{A} contains an identity matrix in itself.
- ▶ For each i , $f(i)^{\text{th}}$ row of \mathbf{M} is just a copy of $f(i)^{\text{th}}$ row of \mathbf{W} .
- ▶ If anchor rows are identified, \mathbf{W} can be found exactly.
 - ▶ Hence we focus of anchor row detection in this presentation.
- ▶ Other rows of \mathbf{A} can be identified easily using a linear program (under some assumption).



Picture Courtesy : A. Moitra, "Computing a Nonnegative Matrix Factorization – Provably," Dec. 2011.

Geometric Picture of the Problem

Definition

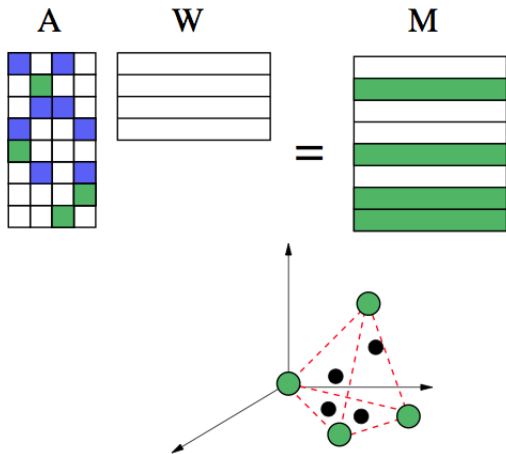
A nonnegative matrix \mathbf{W} is simplicial if no row in \mathbf{W} can be expressed as a linear convex combination of the remaining rows.

Lemma

If a nonnegative matrix \mathbf{M} has a separable factorization \mathbf{AW} , then there is one in which \mathbf{W} is simplicial.

- ▶ This implies that anchor rows correspond to the extreme points of the convex hull containing other rows.

Geometric Picture of the Problem (cont.)



Picture Courtesy : A. Moitra, "Computing a Nonnegative Matrix Factorization – Provably," Dec. 2011.

Anchor Detection Algorithm (Noise Free Case)

Definition

The row \mathbf{M}_j is a *loner* if (ignoring other rows that are its copies), it is not in the convex hull of the remaining rows.

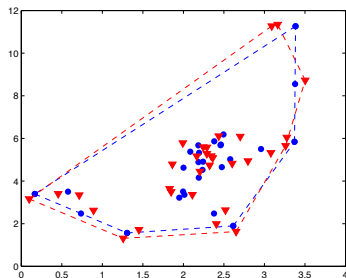
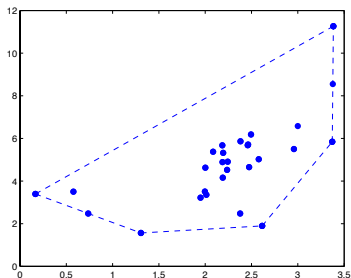
- ▶ Using Linear Programming, it is easy check whether a point is a loner.

Lemma

A row is a loner iff it is an anchor row (equal to some row \mathbf{W}_j).

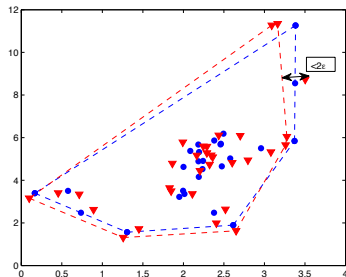
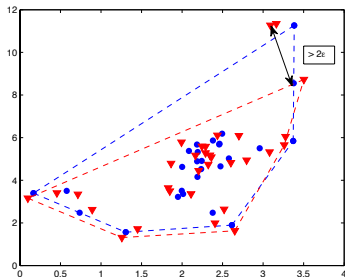
Anchor Detection Algorithm (Small Perturbation Case)

- ▶ Why can't we just apply the previous algorithm?
 - ▶ There may be multiple copies of an extreme point.
 - ▶ Their distance to the convex hull the remaining rows may be indistinguishable with these distances for the ones on/near some face of the simplex.



Small Perturbation Case (cont.)

- ▶ Main idea : remove all points in an small ℓ_1 neighborhood of the point which is going to be tested, then check whether it is loner.
 - ▶ What should be the size of that neighborhood in order to guarantee a robust anchor row detection?
 - ▶ What is the right threshold for the distance of a loner to the convex hull of remaining rows?
 - ▶ How close would be the robust loners to the rows of \mathbf{W} ?



Small Perturbation Case (cont.)

Definition

\mathbf{W} is called α -robust simplicial, if ℓ_1 distance of any row in \mathbf{W} to the convex hull of the remaining rows is larger than α .

- ▶ This condition holds (with high probability in terms of n) for many reasonable prior distributions on \mathbf{W} .

Definition

A point \mathbf{M}'_j is called a *robust loner*, if ignoring rows that are within ℓ_1 distances of $d = 10\epsilon/\alpha + 2\epsilon$, its ℓ_1 distance to convex hull of the remaining rows be more than 2ϵ .

- ▶ Robust loners can be found using Linear Programming.

Main Facts about Robust Loners

Lemma

For all i , $\mathbf{M}'_{f(i)}$ will be a robust loner.

- ▶ Remind that $\|\mathbf{W}_i - \mathbf{M}'_{f(i)}\|_1 \leq \epsilon$.
- ▶ Hence for this set of robust loners, the retrieved row is within $\epsilon \ell_1$ distance of the ground truth.

Lemma

If \mathbf{M}'_j has ℓ_1 distance more than $d + \epsilon = 10\epsilon/\alpha + 3\epsilon$ to all \mathbf{W}_i 's, then it can't be a robust loner.

- ▶ These two lemmas imply that
 - ▶ For all \mathbf{W}_i 's, we retrieve at least one row of \mathbf{M}' which is ϵ close to it.
 - ▶ The other retrieved rows have distance less than $\mathcal{O}(\epsilon/\alpha)$ to some \mathbf{W}_j .
 - ▶ ϵ should be small with respect to α to have reasonable estimation of \mathbf{W} .

Concluding Remarks

- ▶ Using row normalization + assuming simplicial second factor + Separability assumption implies the possibility of totally efficient NMF in the noise free case.
- ▶ Using same components as the previous one and assuming robust simplicial factor implies possibility of totally efficient approximate NMF in the small perturbation case.